

FIG. 1A

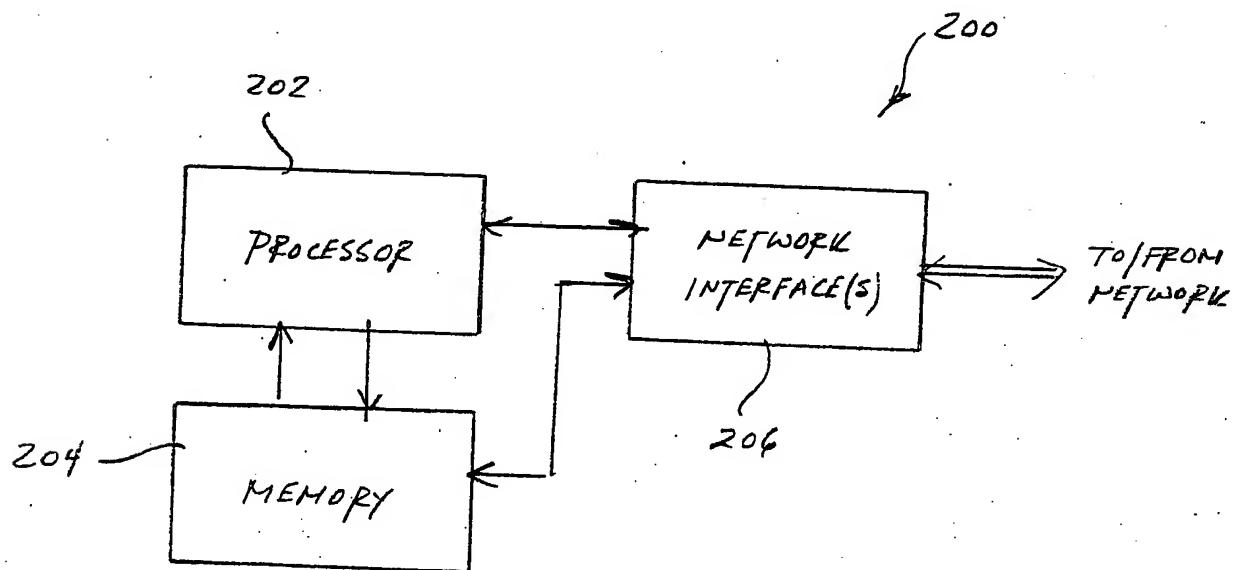


FIG. 1B

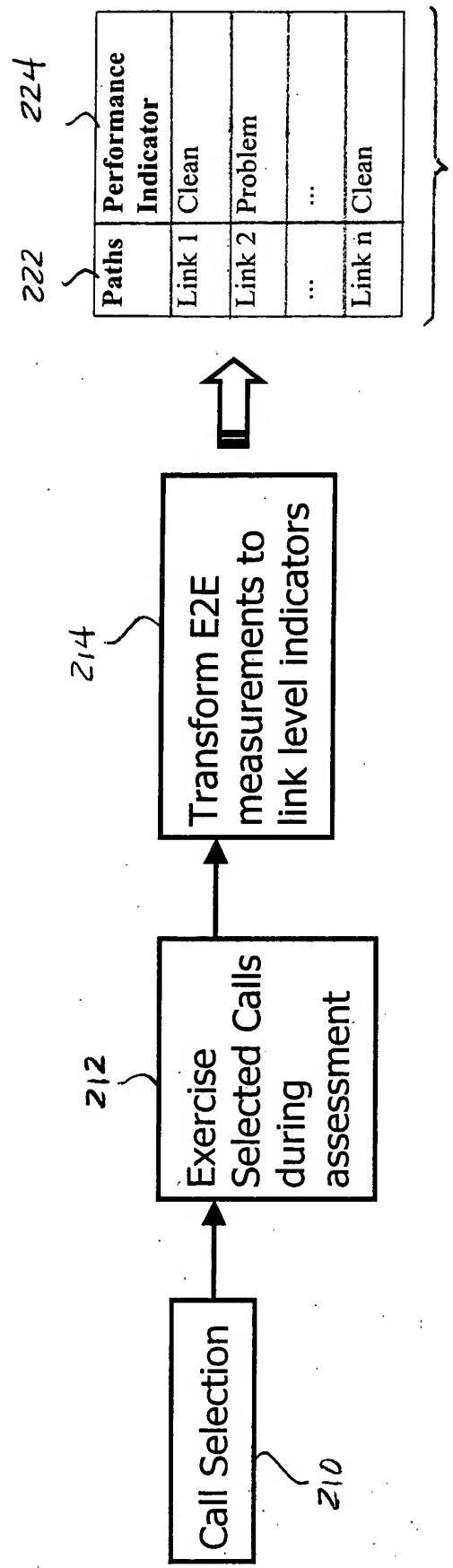


FIG. 2A

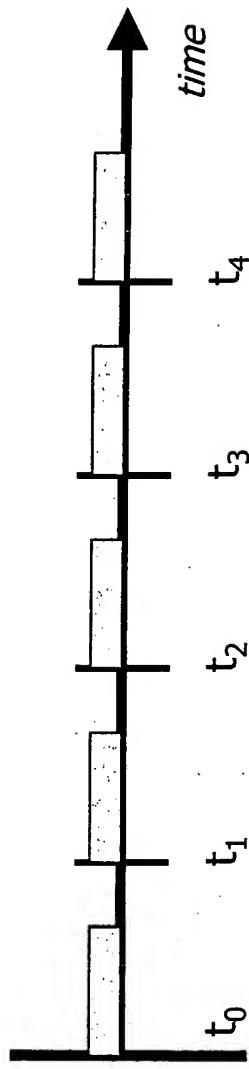
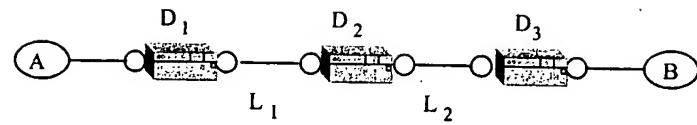
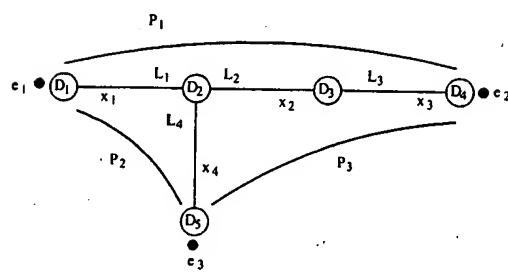


FIG. 2B

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FIG, 3



FIG, 4

	$L_1$	$L_2$	$L_3$	$L_4$
$P_1$	1	1	1	0
$P_2$	1	0	0	1

Flow matrix 1

	$L_1$	$L_2$	$L_3$	$L_4$
$P_1$	1	1	1	0
$P_2$	1	0	0	1
$P_3$	0	1	1	1

Flow matrix2

FIG, 5

Equations with Flow matrix 1

$$\begin{aligned}x_1 + x_2 + x_3 &= y_1 \\x_1 + x_4 &= y_2\end{aligned}$$
$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ \hline x_1 & x_2 & x_3 & x_4 \end{array} \right] = \left[ \begin{array}{c} y_1 \\ y_2 \\ \hline y_1 \\ y_2 \end{array} \right]$$

Equations with Flow matrix 2

$$\begin{aligned}x_1 + x_2 + x_3 &= y_1 \\x_1 + x_4 &= y_2 \\x_2 + x_3 + x_4 &= y_3\end{aligned}$$
$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ \hline x_1 & x_2 & x_3 & x_4 \end{array} \right] = \left[ \begin{array}{c} y_1 \\ y_2 \\ y_3 \\ \hline y_1 \\ y_2 \\ y_3 \end{array} \right]$$

FIG. 6

Generate\_Pipes( $G = (D, L)$ ):Network Topology Graph,  $E$ : Set of Leaves)

$I$ : Set of pipes in  $G$  wrt  $E$

$I \leftarrow \emptyset$

Compute  $P$  for  $G$  wrt  $E$

Let  $M$  be the complete flow matrix for  $G$  and  $P$

// Group links with the same column vector into disjoint sets

Let  $k$  be the number of distinct column vectors in  $M$

Form a set  $\mathbf{S} = \{S_0, S_1, \dots, S_k\}$  where :

each  $S_i$ ,  $0 < i \leq k$  contains links in  $L$  with the  $i^{th}$  distinct column vector in  $M$

// Ensure that links in each element of  $\mathbf{S}$  form a path in  $G$

for  $i = 1$  to  $|\mathbf{S}|$

if links in  $S_i$  are consecutive and form a path

then merge  $S_i$  into path p,  $I \leftarrow I \cup \{p\}$

else  $I \leftarrow I \cup S_i$

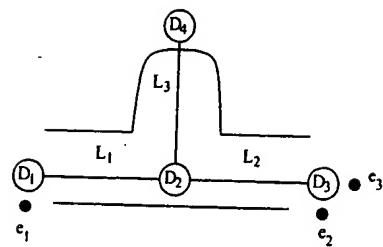
return  $I$

// add the path formed by the links in  $S_i$  as a pipe  
// add each link as a pipe by itself

FIG. 7

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$$\left[ \begin{array}{c|ccc} & L_1 & L_2 & L_3 \\ \hline e_1 - e_2 & 1 & 1 & 0 \\ e_1 - e_3 & 1 & 1 & 2 \\ e_2 - e_3 & 0 & 0 & 0 \end{array} \right]$$

FIG. 8

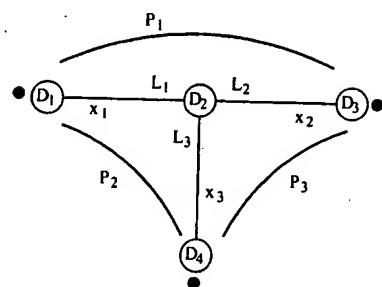


FIG. 9

Flow matrix 1

$$\left[ \begin{array}{ccc} 1 & 1 & 0 \\ 1 & 0 & 1 \end{array} \right]$$

Flow matrix 2

$$\left[ \begin{array}{ccc} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{array} \right]$$

FIG. 10

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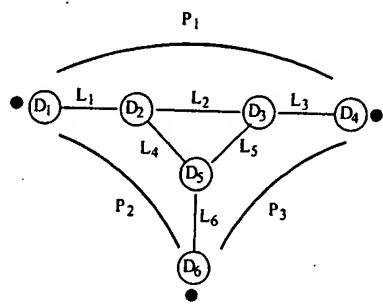


FIG. 11

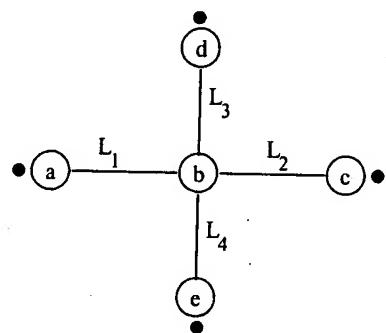


FIG. 12

	$L_1$	$L_2$	$L_3$	$L_4$
$L_1.L_4$	1	0	0	1
$L_4.L_2$	0	1	0	1
$L_2.L_3$	0	1	1	0
$L_3.L_1$	1	0	1	0

$$\{L_1.L_4, L_4.L_2, L_2.L_3, L_3.L_1\}$$

	$L_1$	$L_2$	$L_3$	$L_4$
$L_1.L_2$	1	1	0	0
$L_1.L_3$	1	0	1	0
$L_1.L_4$	1	0	0	1
$L_2.L_3$	0	1	1	0

$$\{L_1, L_2, L_3, L_4\}$$

FIG. 13

**Select\_Matrix( $G' = (D', I)$ :Reduced Network Topology Graph,  $E$ : Set of Leaves)**

```

 $W$ : Set of worms in  $G'$  wrt  $E$ ,  $W \leftarrow \emptyset$ 
 $R$ : Set of paths,  $R \leftarrow \emptyset$ 
Compute  $P'$  for  $G'$  wrt  $E$ 
 $open \leftarrow P'$ 
while  $open \neq \emptyset$ 
    select  $p$  from  $open$ 
    for each pipe  $c_i$  on  $p = c_1.c_2 \dots c_{length(p)}$ 
        if  $\exists S \subset open$  such that  $S$  makes  $c_i$  estimable
            Compute  $S'$  which has the original value of each path in  $S$ 
             $R \leftarrow R \cup S'$ 
             $W \leftarrow W \cup \{c_i\}$ 
            update  $open$  and  $W$  such that  $\forall p' \in open$ 
                 $p'$  does not contain any estimable path in  $W$ 
                //  $c_i$  is removed from paths in  $open$ 
        else
             $c_{i+1} \leftarrow c_i.c_{i+1}$ 
             $open \leftarrow open \setminus \{p\}$ 
    return  $W, R$ 

```

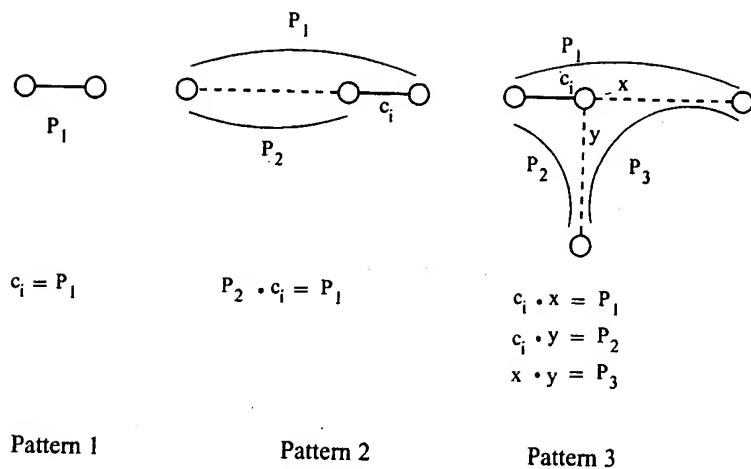


FIG. 15

**Compute\_EstPaths( $G' = (D', I)$ :Reduced Network Topology Graph,  $E$ : Set of Leaves,  $P'_{t_i}$ : End-to-end paths at time  $t_i$ )**

```
M: A Minimal set of estimable paths for  $G'$  wrt  $E$ ,  $M \leftarrow \emptyset$ 
open  $\leftarrow P'_{t_i}$ 
while open  $\neq \emptyset$ 
    while open not converged
        select  $p$  from open
        for each pipe  $c_i$  on  $p = c_1.c_2.\dots.c_{length(p)}$ 
            if  $\exists S \subset open$  such that  $S$  makes  $c_i$  estimable
                 $M \leftarrow M \cup \{c_i\}$ 
                update open and  $M$  such that  $\forall p' \in open$ 
                 $p'$  does not contain any estimable path in  $M$  and
                open  $\leftarrow open \setminus \{p\}$ 
            else
                abort processing of  $p$ 
                if open  $\neq \emptyset$ 
                    select shortest  $p$  in open
                    open  $\leftarrow open \setminus \{p\}$ 
                     $M \leftarrow M \cup \{p\}$ 
    return  $M$ 
```

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Fig. 16